CROISSANT: Centralized Relational Interface for Web-scale SPARQL Endpoints

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Linked Data (LD)

- Open data paradigm
- Linking facts in open data



Linked Open Data cloud diagram (2014-08)

• RDF (Resource Description Framework)

• e.g.,	Subject	Predicate	Object
	iiwas:iiwas2017	dc:title	″iiWAS2017″
	iiwas:iiwas2017	dbo:location	dbr:Salzburg
	dbr:Salzburg	dbo:country	dbr:Austria



Search over LD

- Finding facts in LD data
- Standardized method: SPARQL query
 - Graph pattern-based requirement representation
 - Bindings to variables in patterns are results.



Motivation

• LD is

heterogeneous, complicated, and distributed.

- Large number of facts
- Large number of classes (types of entities)
- Complicated multi-graph structure
- Distributed and self-managed data servers
- Users are demanded to describe search intents by SPARQL.
 - Graph patterns
 - → require understanding of the structure of data
 - Multiple queries for distributed data
 require integration of multiple results

Our Objective and Approach

- Objective
 - Queriable interface for popular query languages for distributed RDF data
- Approach: CROISSANT
 - Centralized interface
 unified access interface for distributed data
 - Relational view-based interface
 reducing effort for understanding data
 - Query rewriting from SQL to SPARQL
 - Query optimization

CROISSANT: an overview



<u>Users' point-of-view,</u> <u>SPARQL endpoints are invisible.</u>

CROISSANT: view definition

(name, schema, endpoint_url, SPARQL_query)

- *name*: view name
 - e.g., movie
- schema: relational schema of the view
 - e.g., (movie_id, title, budget)
- endpoint_url: location of SPARQL endpoint
 - e.g., http://dbpedia.org/sparql
- SPARQL_query: SELECT query of data in the view
 - e.g., SELECT ?movie_id ?title ?budget WHERE { ?movie_id rdf:type dbo:Film. ?movie_id dc:title ?title. ?movie_id dbo:budget ?budget. }

CROISSANT: query execution

• Naive execution

- 1. Execute SPARQL query of corresponding views and store the results into local database.
- 2. Perform SQL query over the local database.



Performance issues

- Execution cost of SPARQL queries
 - Immature performance of SPARQL endpoints
- Transportation cost from SPARQL endpoints



Query Optimization

- Basic idea: reduce #results from SPARQL endpoints
- Strategies
 - View materialization
 - Everything is transferred to local database in advance.
 - Suffer from update issues.
 - Projection push-down
 - Selection push-down
 - View query merge

Projection Push-down

• Push projection conditions into view SPARQL queries.

SELECT **title** FROM movie WHERE budget > 100,000

Input SQL

SELECT <u>?movie_id</u> **?title** ?budget WHERE { ?movie_id rdf:type dbo:Film. ?movie_id dc:title ?title. ?movie_id dbo:budget ?budget. }

Pushed-down View SPARQL query

Selection Push-down

- Push selection conditions into view SPARQL queries.
- Pushing rules

Comparator	Variable type	FILTER expression
θ	Numeric	FILTER (variable θ value)
=	Textual	<pre>FILTER (str(variable) = value)</pre>
like	Textual	<pre>FILTER regex(variable, value)</pre>

SELECT title FROM movie WHERE **budget > 100,000**

Input SQL

SELECT ?movie_id ?title ?budget
WHERE { ?movie_id rdf:type dbo:Film.
 ?movie_id dc:title ?title.
 ?movie_id dbo:budget ?budget.
 FILTER (?budget > 100,000) }

Pushed-down View SPARQL query

Naive Join Query Processing

• Steps

1. Materialize joining views

2. Perform relational join



View Query Merge – motivation

- Naive join is not efficient if two views are of same SPARQL endpoints.
 - Two separate SPARQL queries are performed.
 - Two sets of larger results are transferred.



Results

Local database

View Query Merge – approach

• Idea:

Combine two view SPARQL query on the same SPARQL endpoint into a single query.

- Approach:
 - 1. Combine projection variables into one set.
 - 2. Combine graph patterns into one pattern set.
 - 3. Put FILTER clause for join conditions.
 - If a condition is equality of same attribute names, corresponding FILTER clause is eliminated.

View Query Merge – example

<u>Input SQL</u> SELECT m.title, ma.actor FROM movie m, movie_actor ma WHERE m.movie_id = ma.movie_id

SELECT ?movie_id ?title ?budget
WHERE { ?movie_id rdf:type dbo:Film.
 ?movie_id dc:title ?title.
 ?movie_id dbo:budget ?budget. }

Movie view

SELECT ?movie_id ?actor
WHERE { ?movie_id rdf:type dbo:Film.
 ?movie_id dbo:starring ?actor. }

Movie-actor view

SELECT ?movie_id ?title ?budget ?actor
WHERE { ?movie_id rdf:type dbo:Film.
 ?movie_id dc:title ?title.
 ?movie_id dbo:budget ?budget.
 ?movie_id rdf:type dbo:Film.
 ?movie_id dbo:starring ?actor.
 FILTER (movie_id = movie_id).
}

Merged SPARQL query

Experimental Evaluation

- Objective: Check efficiency of CROISSANT
 Efficiency: query execution time
- Dataset: DBpedia (http://dbpedia.org/sparql)
- Views: movie, actor, movie_actor
- Queries (on the same SPARQL endpoint)
 - Two selection queries of different selectivity
 - Observe effect of selectivity
 - Two join queries w/ and w/o selection condition
 - Observe effect of view query merge
 - Observe overall performance of CROISSANT

Experimental Results

Bars P: Pure execution M: Materialization SP: Selection push-down MG: View query marge MGSP: SP + MP

<u>Insights</u>

- Materialization is obviously the best.
- Selection push-down works when low selectivity.
- View query merge is powerful for join queries.



Conclusion & Future Work

• CROISSANT

- Centralized interface
- Relational view-based interface
- Query rewriting from SQL to SPARQL
- Query optimization
 - Experimental evaluation introduces optimization techniques work.
- Future work
 - Update issue for taking full advantage of materialization
 - Unified optimization with materialization